

#### CPU Critical: A Modern Revisit of a Classic WLM Option



z/OS Performance Education, Software, and Managed Service Providers



Creators of  $\mathsf{Pivotor} \mathbb{R}$ 

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#### Abstract



#### WLM and CPU Critical Control

Buckle up, folks! As of z/OS 3.1, all WLM importance level 1 service class periods will be treated as CPU critical. Why buckle up? Because this change is one of the rare, incompatible changes made by IBM, and has the potential to negatively affect the WLM management of your environment. Luckily, there is an option to disable this incompatible change. However, there is a parameter to optionally disable this setting, if you feel the incompatible default may negatively impact the way WLM manages your workloads.

During this webinar, Scott Chapman and Peter Enrico will provide a review of the WLM CPU critical control and discuss the changes made to z/OS 3.1 that, by default, causes all importance level 1 service class periods to be treated as CPU critical. They will also feature some Pivotor<sup>®</sup> reports to help you determine whether you will be positively or negatively affected by this change.

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  - WLM and CPU Critical Control
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  - Back to Basics Evaluating Latent Demand
  - Understanding SMF 98 Locking Measurements
  - Overseeing z/OS Performance Management With Your Outsourcer
  - Standard Measurements when Monitoring Transactions
  - Processor Comparison Discussion
  - z/OS Performance Management in an AI World
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### **Presentation Overview**



 z/OS 3.1 is defaulting all importance 1 workloads to "Implicit Long-Term CPU Protection", aka "CPU Critical"

- This will likely lead to dispatching priority changes on most systems.
- Some environments will be significantly impacted

#### This presentation

- Refresher of z/OS CPU Dispatching and Dispatching Priorities
- Refresher of WLM CPU controls and algorithms
- Refresher of CPU Critical Control
- Overview of new Implicit CPU Protection in z/OS 3.1



# Refresher : CPU Dispatching and Dispatching Priorities

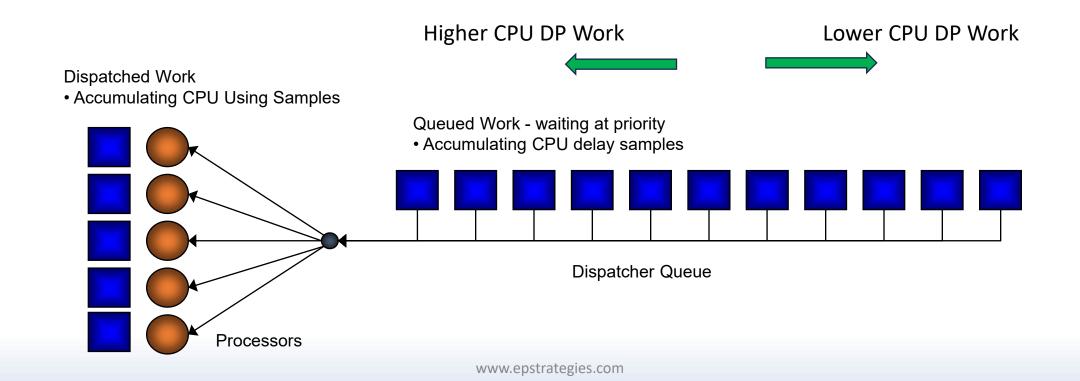
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### **Defining CPU Dispatching Priorities**



• Dispatching priorities determine the order in which operations are executed by the CPU

- All address spaces and enclaves are assigned a CPU dispatching priority
- The dispatching priority order is determined by WLM based on WLM goals
- Work queued to use the CPU is placed onto the dispatching queue in CPU dispatching priority order



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### **CPU Dispatching Priorities Schema**

• The reasons we care about CPU DP order

- Prioritization of CPU access to each workload
- To assist when there are feeder effects

• Workload efficiencies

 Managing goals and processor optimization by addressing processor delays

#### Other reasons include

- CPU promotion for work holding a resource
- Reduced preemption
- Fairshare dispatching

255	SYSTEM
254	SYSSTC
253	
249	Unused
248	Small Consumer
247	
	Priorities Used
	for RT or
	Velocity Periods
	(i.e. lmp 1 – 5)
203	
202	Unused
201	
	Discretionary
	(MTTW)
192	
191	Quiesce

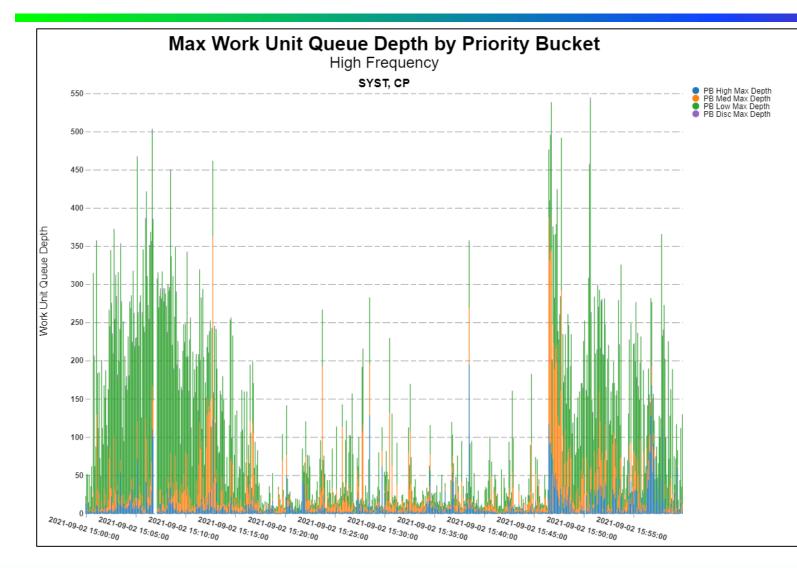




Reason 1: Prioritization of CPU access to each workload

- Example: Higher importance work should run at higher CPU DPs
- Although not a hard and fast rule, to meet goals, higher importance work tends to receive higher CPU DPs from lower importance work
- Relative importance does not translate to relative CPU DPs
  - A higher importance goal could have a lower CPU dispatch priority than a lower importance goal
  - CPU Critical control does influence this
    - Says lower importance work will never have same or higher DP as work identified as CPU critical

### Use SMF 98 to look at dispatch queue depths



Priority bucket statistics

(1=High, 2=Med, 3=Low, 4=Discretionary)

This chart is related to the previous chart, but for CP engines rather than zIIP engines.

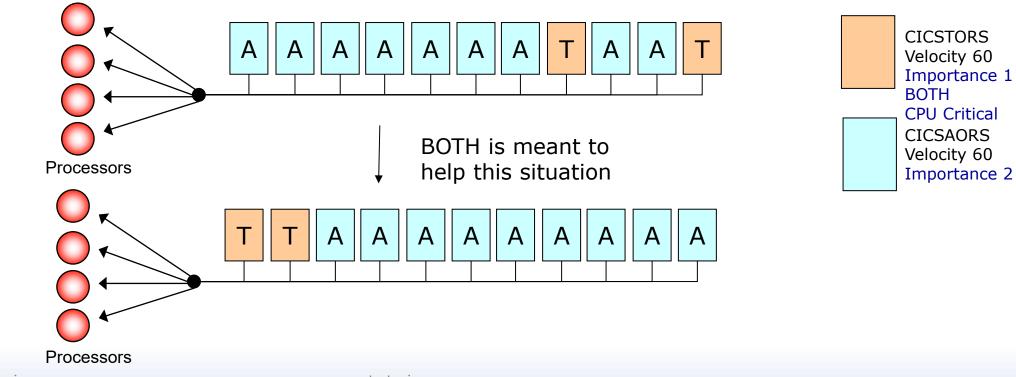
We see that although the dispatching queues are longer. The displaceable work is lower importance. Thus, crossover of higher importance will displace this lower importance work.

### Reason 2: Why we care about CPU DPs

EPS

• Reason 2: To assist when there are feeder effects

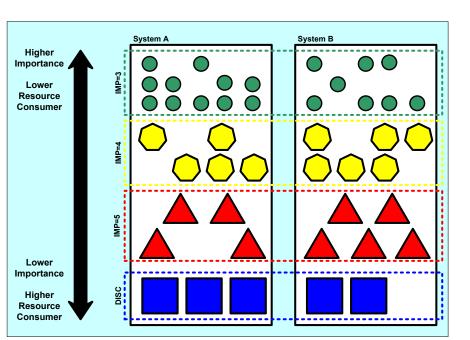
- Example: Certain workloads feed other workloads to progress work
- IRLM -> DB2 (DBM1,MSTR,DIST) -> CICS TOR -> CICS AOR



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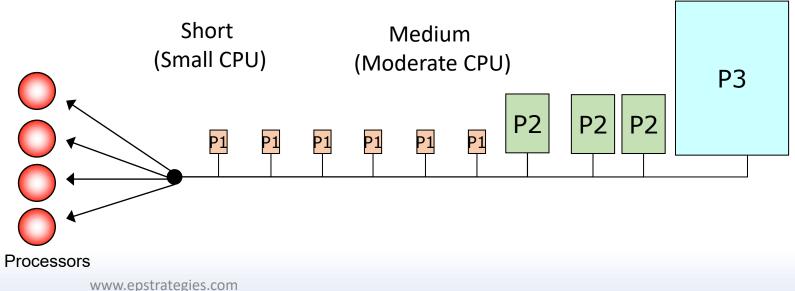
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### Reason 3: Why we care about CPU DPs



- Reason 3: Workload efficiencies
  - Example: Get the shorter work (light CPU consuming work) on the CPU first to get it 'in and out' as quickly as possible
  - This is extraordinarily effective with interactive (DDF, TSO, etc.) and batch workloads

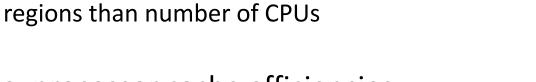
Long (Heavy CPU)



0

- less efficient

- Interactive workloads tend to use processor caches
- Also, processor cache efficiencies



Batch workloads tend to use processor caches more

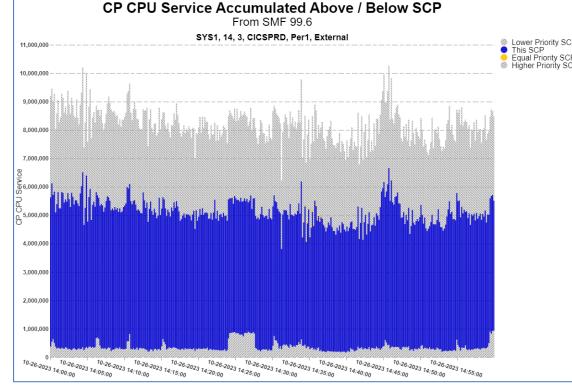
Example: Over initiation of batch, many more CICS

#### Reason 4: Managing goals and optimize processors by addressing processor delays

- - Helps address latent demand
  - 0 Correct CPU dispatching priority helps manage and 0

alleviate latent demand and processor delays





Reason 4: Why we care about CPU DPs

efficiently



• WLM sets dispatching priority for service class periods.

• All address spaces in a service class period have the same base dispatching priority

- Note: there are 'internal periods' and 'external periods'
- There is also something called promotion (discussed later)

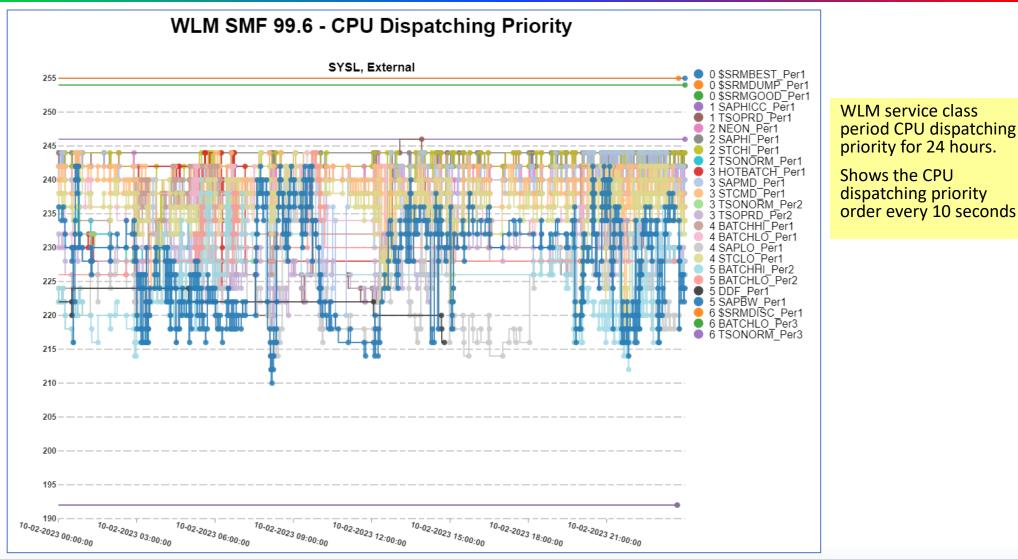
• Multiple service class periods may have the same base dispatching priority.

Unbunching

- There is a concept in the WLM algorithms called 'unbunching'
- After a dispatching priority change, service class periods may be remapped to different dispatching priorities such that there is an unoccupied priority between each occupied priority.

#### SMF 99.6 CPU Dispatching Priority – Every 10 Seconds





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#### CPU Promotion: Another reason for changing CPU DPs:



- WLM briefly 'promotes' individual work units to higher CPU dispatching priorities to boost the access to the CPU for very short periods of time
  - Concept:
    - If a unit of work is regularly demanding CPU and possibly holding a resource, then WLM may briefly
      promote the unit of work to a higher CPU DP to allow it to run for a short period of time in hopes of
      the work unit releasing the resource

#### • Reasons for promotion:

- ENQ Promoted by enqueue management because the work held a resource that other work needed.
- BLK Promoted to help blocked workloads
- LCK Promoted to shorten the lock hold time of a local suspend lock held by the work unit
- SUP Promoted by the z/OS supervisor to a higher dispatching priority than assigned by WLM
- CRM Promoted by chronic resource contention management because the work held a resource that other work needed

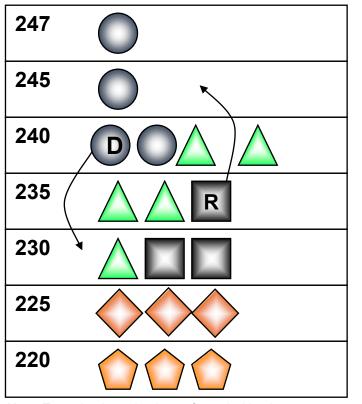


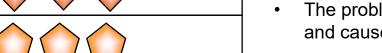
#### WLM CPU Critical Control

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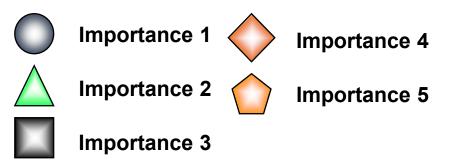
## Overview of WLM CPU Critical Control

• Because some installations are concerned that WLM will not react fast enough for high priority work, WLM has a 'CPU Critical Control'





Note: To make the point, just a few priorities between DP 203 and DP247 are shown.



- With well set predictable goals, DPs tend to be ordered by importance
- If work is missing its goal WLM may decide to adjust its DP equal or above a higher . importance period
- The problem occurs when this lower importance period starts to consume more CPU and causes the higher importance period to miss its goal
- WLM will recognize this condition and fix it, but WLM can be slow to react ٠

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#### CPU Critical aka Long-term CPU Protection



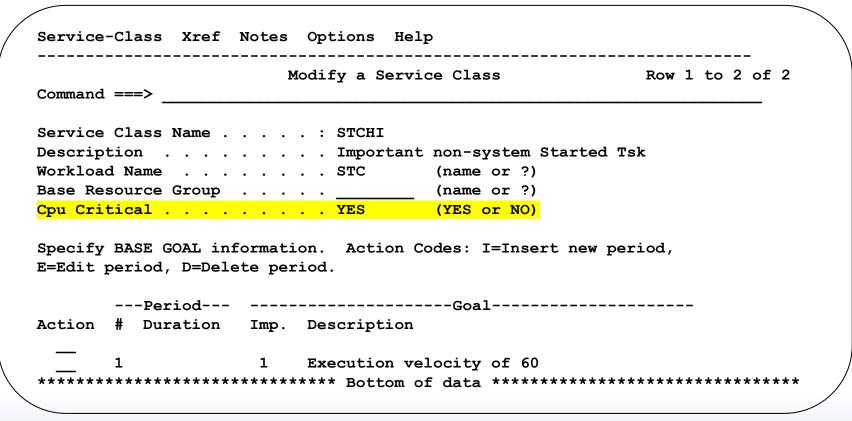
- Long-time option in your WLM service definition
- Enabled by setting YES for CPU Critical on a Service Class
  - Must be a single-period SC and cannot be discretionary
- Ensures that the CPU Critical SC always has a dispatching priority that's greater than the DP of lower importance service class periods
- Note some small amount of lower-importance work may still get higher DP:
  - Due to promotion for locks, resource contention, etc.
  - Small consumers
- General recommendation has been to avoid this option
  - Allows WLM to make better decisions about balancing overall throughput to best meet the goals of all work
- Important: The use of these options limits WLM's ability to manage the system. This may affect system performance and/or reduce the system's overall throughput.
- Note: Implicit CPU Critical for importance 1 work can impact the CPU distribution to lower importance work. Ensure that the goals are appropriate given their importance level. Evaluate the current distribution of CPU at different importance levels, especially those that are covered by CPU Critical to ensure that they have consistent CPU demands.

### Setting CPU Critical Control



• CPU Critical is set at the service class level

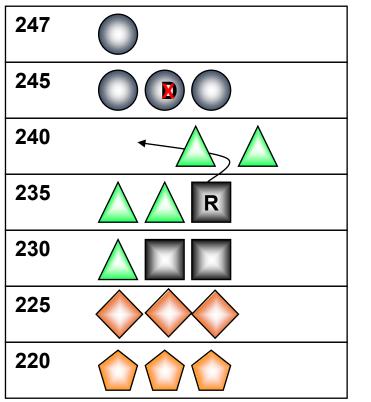
- Can be set for address space, enclave, or transaction-oriented work
- To help ensure that critical work will have a higher CPU DP than lower importance work

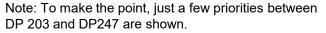


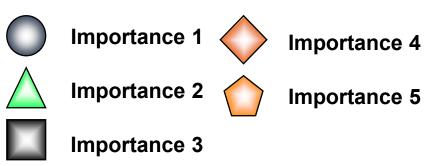
### Effect of WLM CPU Critical Control



 With CPU Critical set for the importance 1 workloads, all lower importance workloads will have a lower dispatching priority







- In the case where that one Importance 1 workload is overperforming and the importance 3 workload is underperforming, now it can't shift those workload DPs around as much to help the importance 3 workload.
- The R workload may be helped somewhat by moving it up to DP 240, but maybe it's not enough because that Importance 1 donor is really consuming a bunch of CPU (and it could suffer some CPU delay and still meet its goal)



#### Implicit CPU Protection in z/OS 3.1

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### z/OS 3.1 Implicit Long Term CPU Protection



New option for "Implicit" Long-Term CPU Protection
 In other words, CPU Critical without having to specify it on every SC definition

#### • Default is "On" for importance 1 service class periods

- Optional, but "Off", for importance 2 and lower service class periods
- only assigned to single-period classes and the first period of multi period service classes.

#### • We think "On" for importance 1 workloads is a bad default

- Could significantly change the dispatching priority of work in the system
- Goes against historical practices of not changing defaults that change behavior

#### • DP/Importance inversions are common

- I.E. Lower Importance work running with a DP above higher importance work
- Not all such inversions are problematic
- Not all importance 1 work really should be importance 1

### IEAOPTxx – CCImp and CCImpBoost



#### • Two new parameters for z/OS 3.1

- CCImp=0|1|2
   (CCImp applies when no boost is active)
- CCImpBoost=0|1|2
   (CCImpBoost applies while the boost is in effect)

#### ° 0

• CPU protection is not assigned to any work unless CPU Critical has been explicitly assigned for the service class in the active WLM policy.

#### ° 1|2

 Specifies the importance level up to which CPU protection is implicitly assigned to the work when CPU Critical has not been assigned for the service class in the active WLM policy.

#### • Examples:

 CImp=0 and CCImpBoost=1 suppresses implicit CPU protection but enables implicit CPU protection when a boost is in effect for importance 1 work.

0

CCImp=1 and CCImpBoost=1 assigns implicit CPU protection for any importance 1 work.

### From the IBM z/OS 3.1 Documentation

(z/OS 3.1 MVS Planning Workload Management SC34-2662)

"Beyond explicitly setting the CPU Critical option for single-period service classes of any importance except discretionary, CPU protection is implicitly assigned for the first period of any service class of importance 1 (and importance 2 when a boost is in effect)."

"If you want to modify the importance level for automatically setting CPU protection, or even disable it, use the **CCImp** and **CCImpBoost** parameters in the **IEAOPTxx** member." (See next slide)

**"Note:** Implicit CPU Critical for importance 1 work can impact the CPU distribution to lower importance work. Ensure that the goals are appropriate given their importance level. Evaluate the current distribution of CPU at different importance levels, especially those that are covered by CPU Critical to ensure that they have consistent CPU demands."

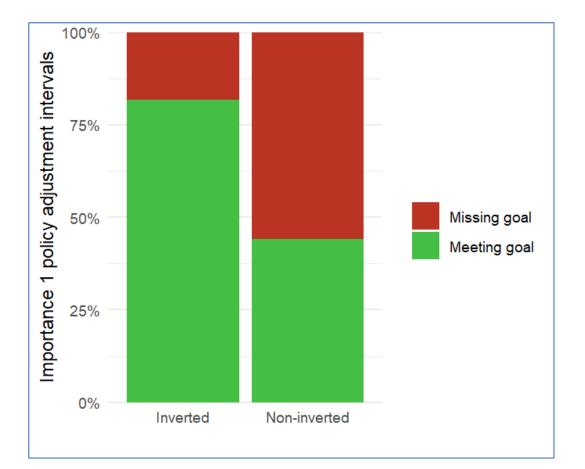
### We did an analysis of 116 systems



- Covered a variety of sizes from large to small, "IPO" to "Prod", a couple of dozen customers
- Evaluated a day's worth of 99.6 data from each system (over 17M records)
- Created 2 new metrics to help understand the risk/benefit:
  - For SCPs that would be bumped down: Inversion Risk Ratio – relative amount of CPU that would move above the SCP
  - For SCPs that would move up in priority: Protection Benefit Ratio – relative amount of CPU that would move below this SCP
  - Higher numbers means more potential risk/benefit
    - Can be very high if there's a relatively large difference in the consumption of the workloads

#### Effectiveness of z/OS pre-3.1 WLM

- z/OS pre-3.1 WLM behavior sufficiently protects importance 1 workloads and properly balances workload performance across all importance levels
- Study shows Importance 1 workloads which are running at a lower priority than lower-importance workloads are already meeting their goals more often than those running above all lower importance work (82% to 44%)
- It is important to note the limitations of this metric, as goals are not always assigned in accordance with actual needs and expectations





### Notes: Inversion Risk Ratio



- The ratio of CPU service of importance 1 work to lower-importance work that is running at a higher dispatching priority (excluding work at DP 248)
  - Basically: more importance 1 work is about to move above a particular SCP, means that SCP is at greater risk
  - E.G. you have SCP A (importance 2) consuming 50 SUs that's running above importance 1 work consuming 5000 SUs. SCP A has an inversion risk ratio of 100.
  - Higher numbers = more risk, but even ratios < 1 represent some risk
  - Risk = Might not get as ready access to CPU, might suffer more CPU delay

### Notes: Protection Benefit Ratio



- The ratio of CPU service of lower-importance work that is running at a higher dispatching priority (excluding work at DP 248) to importance 1 workload
  - Basically: more lower importance work is about to move below an importance 1 SCP, means that importance 1 SCP might get a bigger benefit
  - E.G. you have SCP Z at importance 1 and consuming 80 SUs but it's running below lower importance work that's consuming 4000 SUs. SCP Z has a Protection Benefit Ratio of 50.
  - Higher numbers = more potential benefit, but even ratios < 1 represent some possible benefit.
  - Benefit = might get easter/faster access to CPU, might see less CPU delay

### Histograms



• The next slide has two histograms:

- On the left, the Inversion Risk Ratio across all such service class periods which are running above an importance 1 workload in the 116 systems. (Total ~1 million service class policy adjustment intervals)
- On the right, the Protection Benefit Ratio, across all importance 1 SCPs running below lower-importance workloads in the 116 systems. (Total, ~434,000 service class policy adjustment intervals)

 There are instances of very high ratios in both cases, but overall, there's more SCPs at risk than they are that stand to benefit.

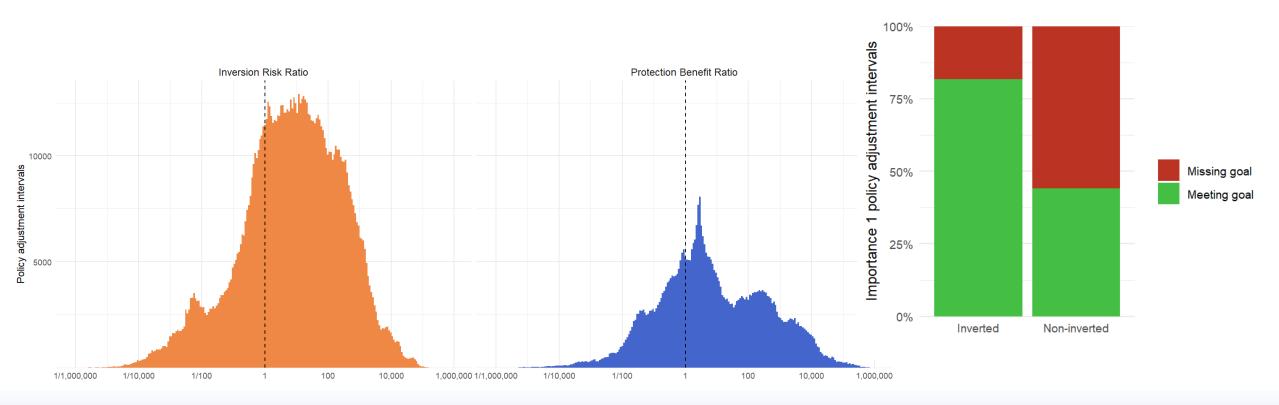
• And most of SCPs that might benefit are already meeting their goals

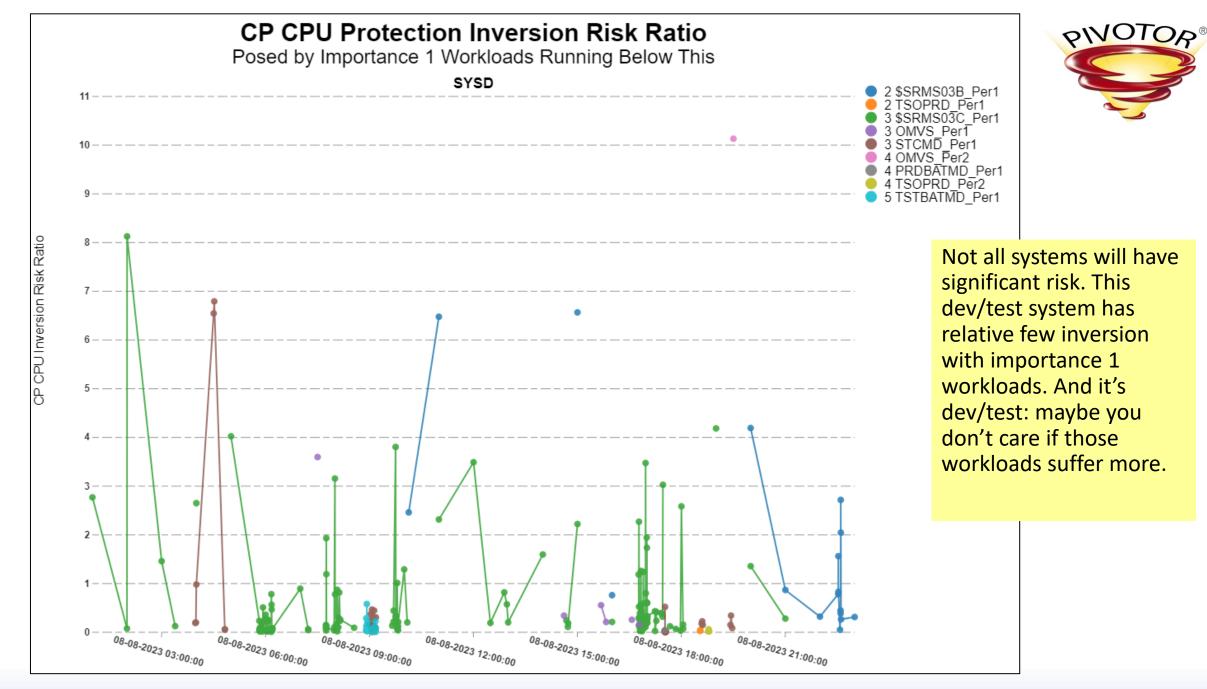
 For those that have set goals appropriate to the workload's actual performance needs, the risks may very well outweigh any potential benefits

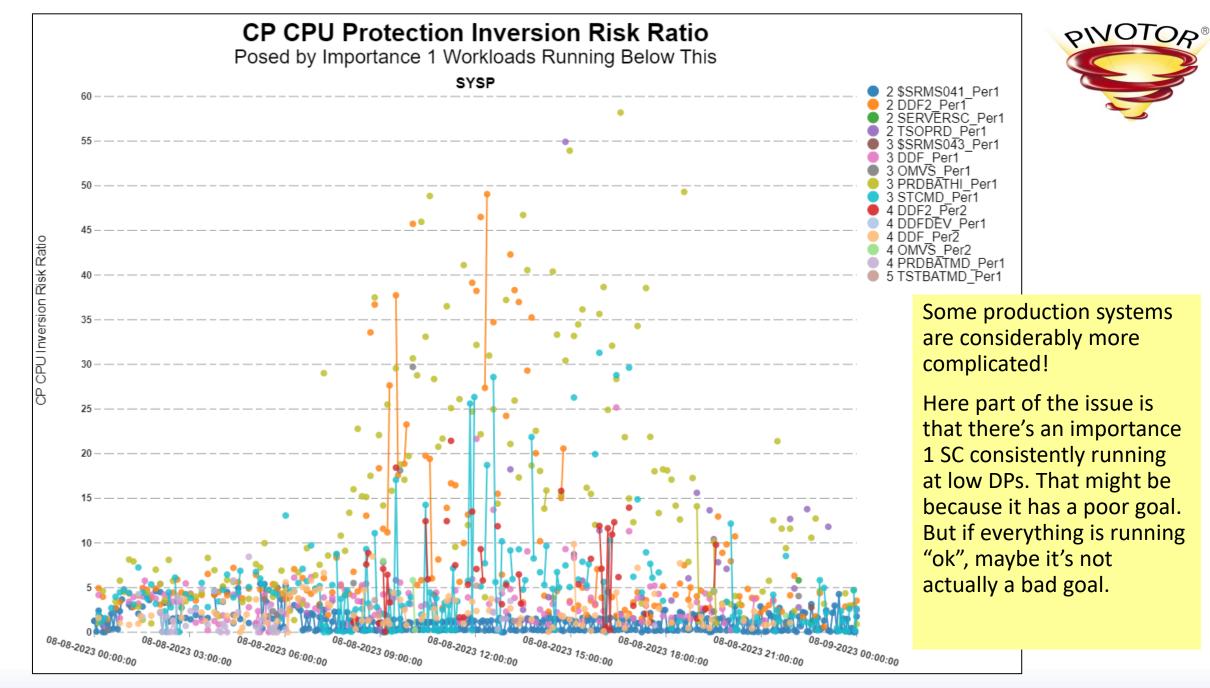
## Study findings



- 78% of systems had at least one interval with an inversion
- 39% of systems had inversions in at least 25% of their intervals
- 82% of "Inverted" Importance 1 workloads were meeting their goal







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### The dumbing down of WLM



In the age of AI, it is strange that IBM is choosing to dumb down WLM

- Customers could already choose to do this... so why the new forced default?
- Also, the disablement is in the IEAOPTxx rather than the WLM service definition
- Defaulting all importance 1 workloads to CPU Critical may have significant impacts
  - "Fixing" these inversions is likely to have a larger negative impact on the lower importance workload than it will a positive impact on the importance 1 workloads
- Making this the new default disregards:
  - The practice of avoiding changing defaults that would change the behavior of the system
  - Long-standing recommendations to generally avoid the use of CPU Critical

 It is not correct to assume that all work is classified to the proper importance and given a proper goal that is consistent with the business and technical requirements, but it's also not correct to assume goals were improperly set as well.

### Our thoughts



- We don't see the need for this change
  - A significant part of the premise of WLM was that it would manage dispatching priorities and could intelligently move them in possibly counter-intuitive ways to better balance throughput for diverse workloads
  - If you want, you can make all importance 1 work CPU Critical today (at least for single-period SCs)
- We'd recommend turning this off for z/OS 3.1 and wish that was the default
- If you want to go to z/OS 3.1 with it on, we might suggest
  - **1.** Evaluate which workloads are at risk
  - 2. Before z/OS 3.1, incrementally add CPU Critical to importance 1 workloads
    - If something goes wrong, you can back out your change and z/OS 3.1 doesn't get the blame
- We do sometimes recommend CPU Critical, but it's an exception, not the rule
- Again: emerging area of study, we might refine our recommendations over time



• White paper available:

Summary of Potential Impact of Implicit Long-Term CPU Protection

• For a copy, please send an email to <u>performance.questions@epstrategies.com</u>

#### • Special thanks to Ethan Chapman for his statistical and R expertise!

#### • QUESTIONS?